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We need better C tools

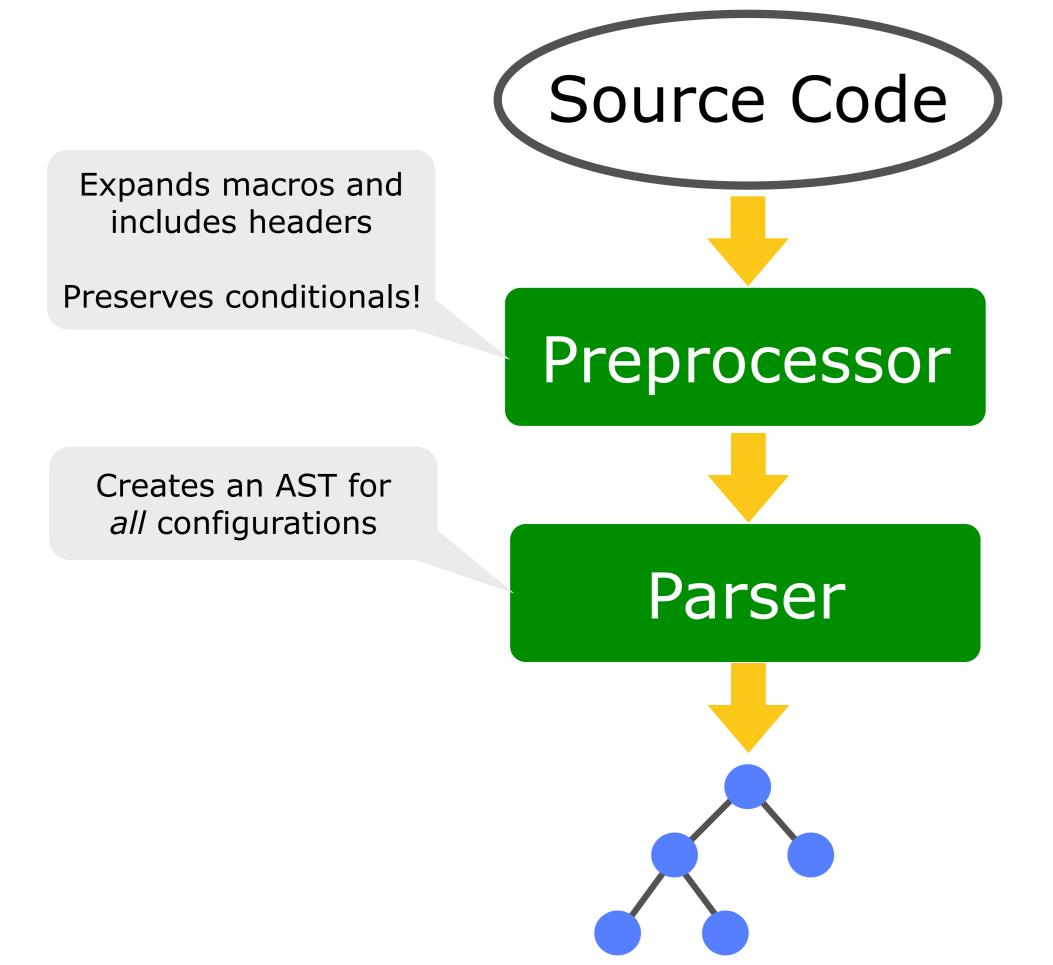
- Linux x86 is large and complex
- Need source code browsers
- 7,500+ compilation units, 5.5 million lines
- Need bug finders
- 1,000 found by static checkers [Chou et al., SOSP '01] • Need refactoring tools
- 150+ errors due to interface changes [Padioleau et al., EuroSys '08] • These tools all need to parse C first

C source code written in both C and the preprocessor

• Source code contains many programs

- #ifdef CONFIG_USB_DEVICEFS
- extern int usbfs_init(void); #else
- static inline int usbfs_init(void){return 0;} #endif
- Linux x86 has 6,000 configuration variables, 2^{6,000} combinations
- Turning on all configuration variables yields only 80% of code [Tartler et al., OSR '11]
- Macros expand to arbitrary C fragments #define for_each_class(c) \ for (c = highest_class; c; c = c->next)
- Directives appear between arbitrary C fragments #define for_each_class(c) \ for (c = highest_class; c; c = c->next)

Solution Approach



SuperC: Parsing All of C by Taming the Preprocessor

How SuperC Works

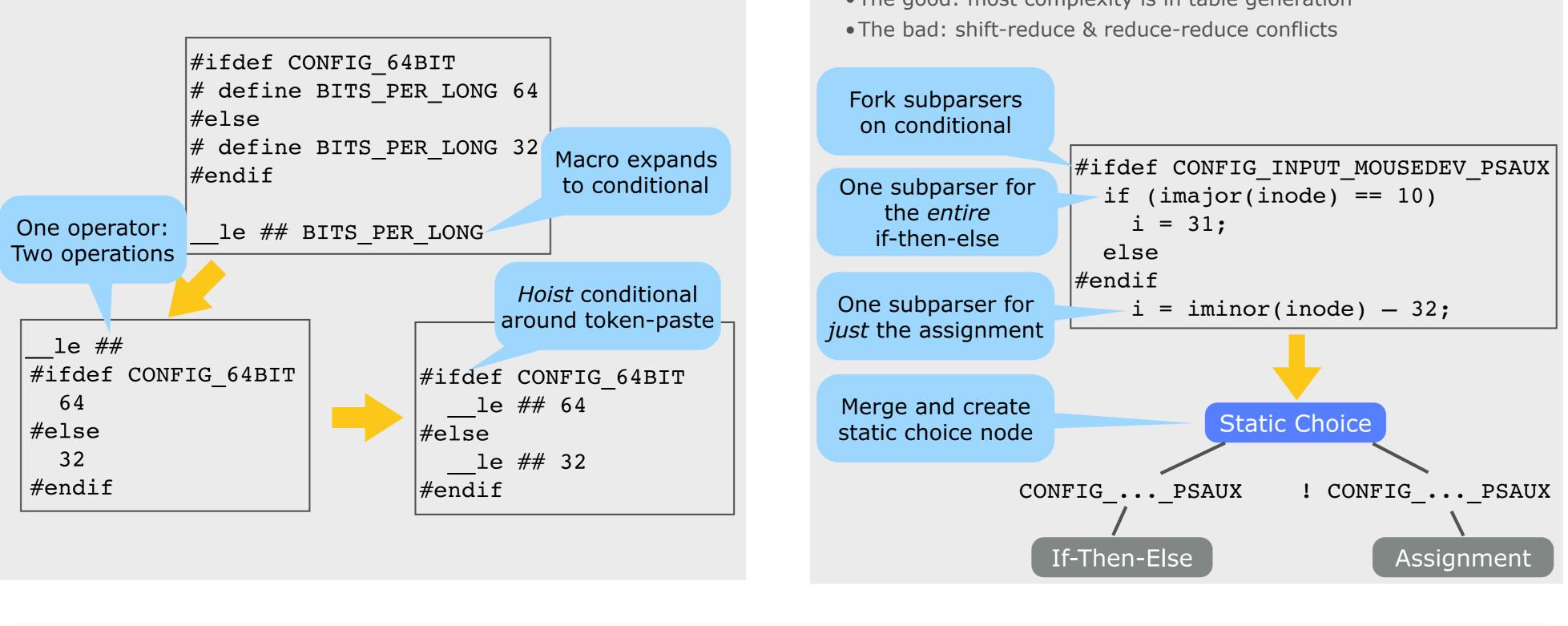
The Preprocessor

Conditionals Invade the Preprocessor

- The preprocessor leaves condtiionals in place
- Conditionals then compose with most preprocessor operations
- Many operations require *hoisting*

The Power of Hoisting

- Works on: token-pasting, stringification, includes, conditional expressions, macros
- Iterates over conditional branches
- Recurses into nested conditionals
- Duplicates tokens across inner-most branches



Evaluation



http://cs.nyu.edu/xtc



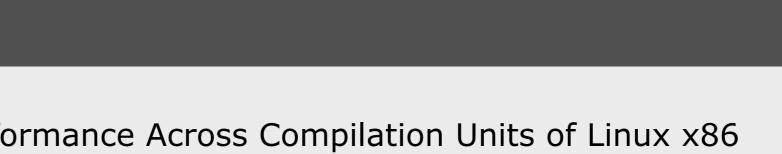
Parsing All Configurations

• Forks subparsers at conditionals • Merges subparsers in the same state after conditionals • Joins AST subtrees with static choice nodes • Preserves mutually exclusive configurations

History Repeats Itself: LR Subparsers

Organizes state in stacks

- Easy forking and merging with DAG
- Is table-driven
- Good performance
- Reuses existing tools and grammars
- The good: most complexity is in table generation

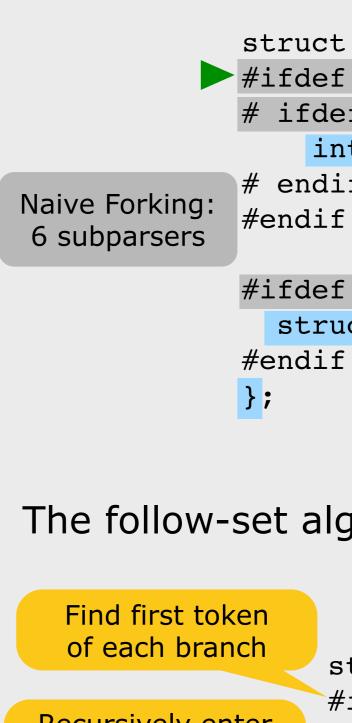


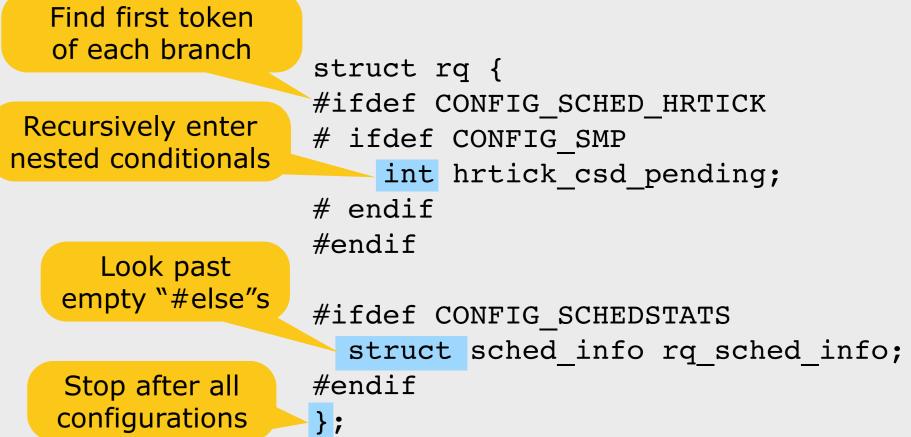
When to Fork Subparsers?

- Blows up on Linux x86

- Across all configurations

Follow-set forks fewer subparsers





Follow-set supports further optimization

- Shared reductions
- Limits redundant work by subparsers
- Lazy shifts
- Limits number of subparsers needed
- Early Reduces
- Improves chances of merging

Parsing Real-World C

• Naive strategy: fork on every conditional branch • Conditionals are 40 levels deep, 10 in a row • Our forking strategy: token follow-set • All tokens reachable from current position

cuct rq {	
def CONFIG_SCHED_HRTICK	
fdef CONFIG_SMP	
<pre>int hrtick_csd_pending</pre>	;
endif ndif	Follow-set: 3 subparsers
def CONFIG_SCHEDSTATS	
<pre>struct sched_info rq_sched_info;</pre>	
2; E	

The follow-set algorithm in action

• Reduce one stack for many follow-set tokens before forking

• Only fork tokens in the nearest condtiional

• Pick reducing subparser before a shifting one